



SNS COLLEGE OF TECHNOLOGY

Coimbatore-35
An Autonomous Institution



Accredited by NBA – AICTE and Accredited by NAAC – UGC with 'A++' Grade
Approved by AICTE, New Delhi & Affiliated to Anna University, Chennai

DEPARTMENT OF ECE

23ECB222- Digital Principles and Computer Organization

II YEAR/ III SEMESTER

1

UNIT 1 – MINIMIZATION TECHNIQUES AND LOGIC GATES

TOPIC - BOOLEAN EXPRESSIONS, MINIMIZATION OF BOOLEAN EXPRESSION



MINIMIZATION OF BOOLEAN ALGEBRA



What is Minimization?

- A Boolean expression is composed of variables and terms. The simplification of Boolean expressions can lead to more effective computer programs, algorithms and circuits.



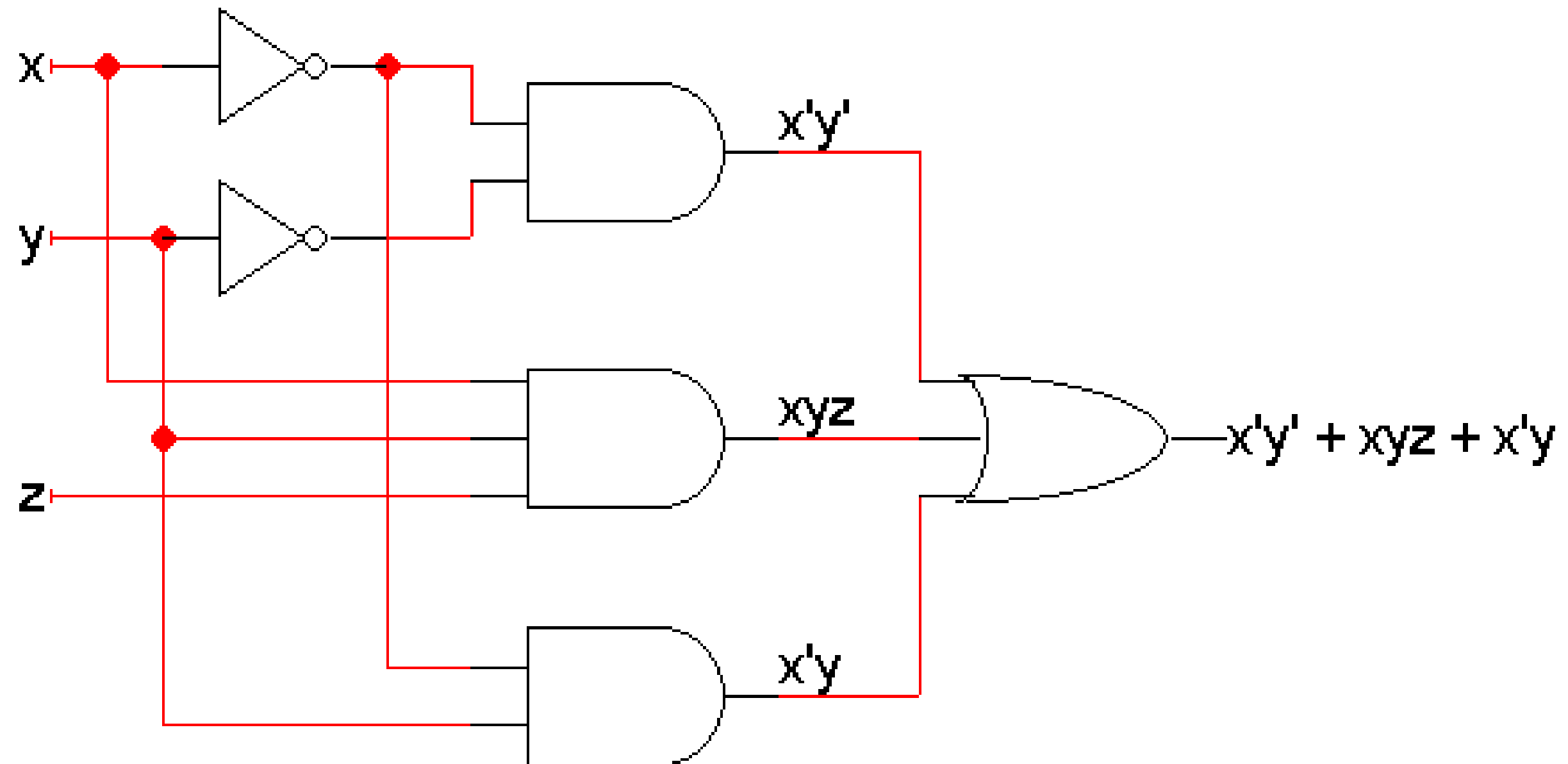
MINIMIZATION METHODS



- Minimization can be achieved by a number of methods, three well known methods are:
 1. Algebraic Manipulation of Boolean Expressions
 2. Tabular Method of Minimization
 3. Karnaugh Maps



Algebraic Manipulation of Boolean Expressions

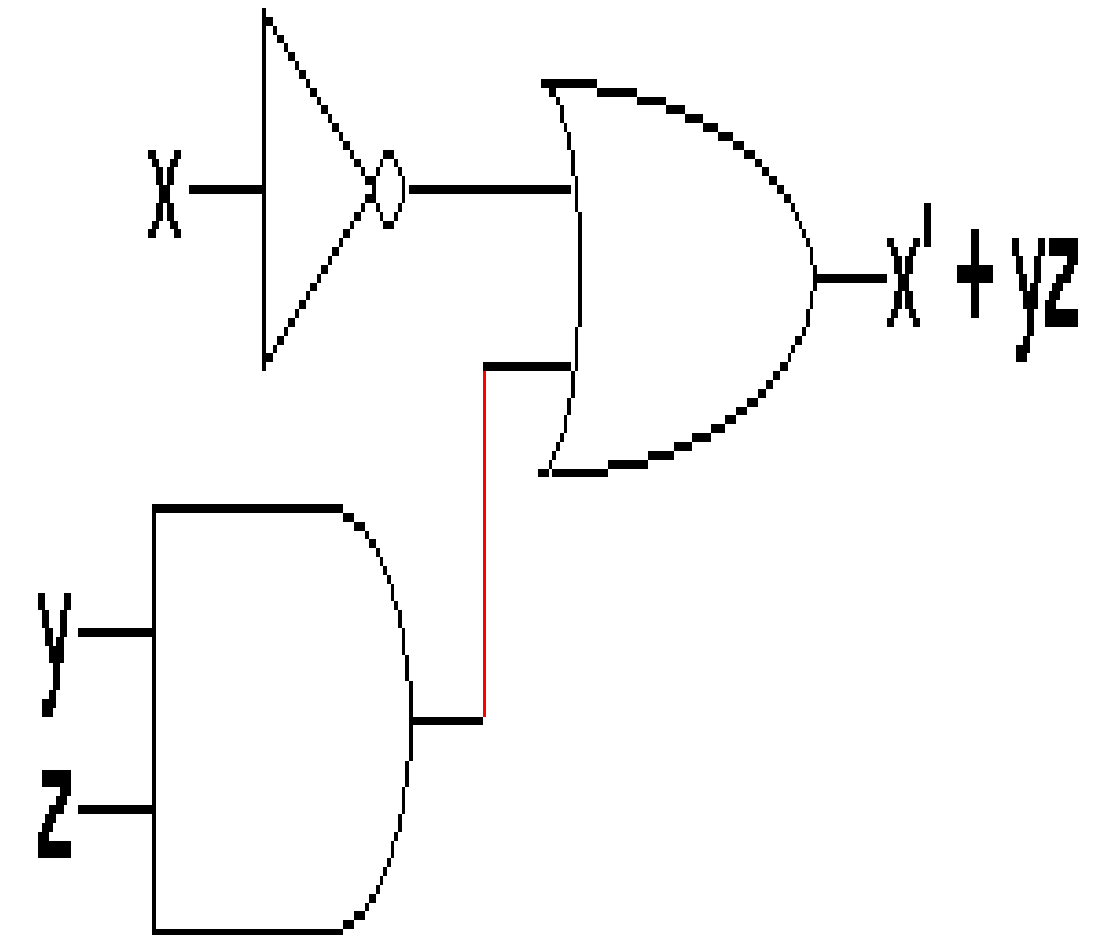




Algebraic Manipulation of Boolean Expressions



- Here are two different but *equivalent* circuits.
- In general the one with fewer gates is “better”:
 - It costs less to build
 - It requires less power
 - But we had to do some work to find the second form





ALGEBRAIC MANIPULATION



EXAMPLE 1

$$\begin{aligned} & x' y' + xyz + x' y \\ &= x' (y' + y) + xyz \quad [\text{Distributive: } x' y' + x' y = x' (y' + y)] \\ &= x' \cdot 1 + xyz \quad [\text{complement: } x' + x = 1] \\ &= x' + xyz \quad [\text{identity: } x' \cdot 1 = x'] \\ &= (x' + x)(x' + yz) \quad [\text{Distributive}] \\ &= 1 \cdot (x' + yz) \quad [\text{complement: } x' + x = 1] \\ &= x' + yz \quad [\text{identity}] \end{aligned}$$



PROBLEMS-BOOLEAN MINIMIZATION

- $AB + \bar{A}C + BC = AB + \bar{A}C$ (Consensus Theorem)

Proof Steps

$$\begin{aligned} & AB + \bar{A}C + BC \\ = & AB + \bar{A}C + 1 \cdot BC \\ = & AB + \bar{A}C + (A + \bar{A}) \cdot BC \\ = & AB + \bar{A}C + ABC + \bar{A}BC \\ = & AB + ABC + \bar{A}C + \bar{A}CB \\ = & AB \cdot 1 + ABC + \bar{A}C \cdot 1 + \bar{A}CB \\ = & AB(1+C) + \bar{A}C(1+B) \\ = & AB \cdot 1 + \bar{A}C \cdot 1 \\ = & AB + \bar{A}C \end{aligned}$$

Justification

Identity element

Complement

Distributive

Commutative

Identity element

Distributive

$1+X = 1$

Identity element



◆ Example 1: A two-level logic expression

$$Z = A'BC + AB'C' + AB'C + ABC' + ABC$$

$$= AB'C + AB'C' + A'BC + ABC' + ABC$$

$$= AB'(C + C') + A'BC + AB(C' + C)$$

$$= AB' + A'BC + AB$$

$$= AB' + AB + A'BC$$

$$= A(B' + B) + A'BC$$

$$= A + A'BC$$

rearrange
distributive
comp.
rearrange
distributive
comp.

◆ Use absorption #2D $\{(X \cdot Y) + Y = X + Y\}$ with $X=BC$ and $Y=A$

$$Z = A + BC$$



EXAMPLE



- $(A + B)(A + C) = A + BC$
- This rule can be proved as follows:
- $(A + B)(A + C) = AA + AC + AB + BC$ (Distributive law)
 $= A + AC + AB + BC$ ($AA = A$)
 $= A(1 + C) + AB + BC$ ($1 + C = 1$)
 $= A \cdot 1 + AB + BC$
 $= A(1 + B) + BC$ ($1 + B = 1$)
 $= A \cdot 1 + BC$ ($A \cdot 1 = A$)
 $= A + BC$



ASSESSMENT TIME



SOLVE THE EXPRESSIONS USING BOOLEAN LAWS

$$1. F(A, B, C) = A'B + BC' + BC + AB'C'$$

$$2. F(A, B, C) = (A+B)(A+C)$$



REFERENCES



- 1.M. Morris Mano, “Digital Design” 4TH Edition PHI/2008, Singapore Pvt.Ltd,new Delhi 2003.
- 2.John.M Yarbrough, “Digital Logic Applications and Design”, Thomson Learning, 2006.



THANK YOU